

Accelerating Numerical Software Libraries with Multi-Precision Algorithms

IDEAS Webinar

Wednesday, May 13, 2020

Hartwig Anzt and Piotr Luszczek



IDEAS
productivity



ECP
EXASCALE COMPUTING PROJECT



better
scientific
software

Who are we?

- Hartwig Anzt
 - Mathematical library developer;
 - Lead of the ECP Multiprecision effort;



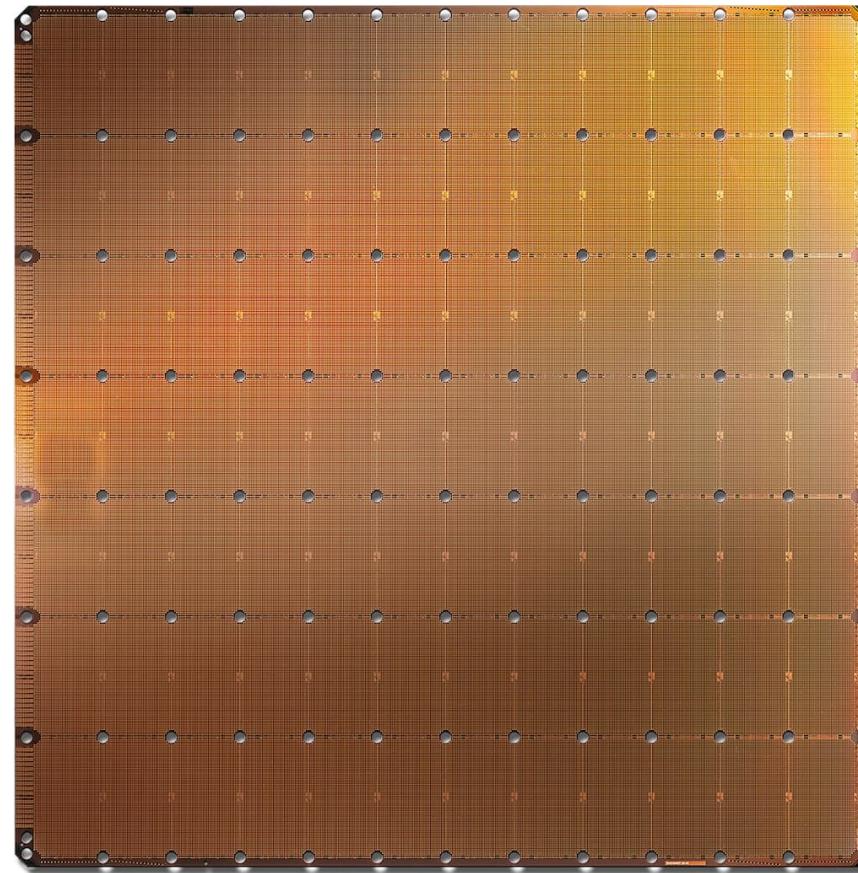
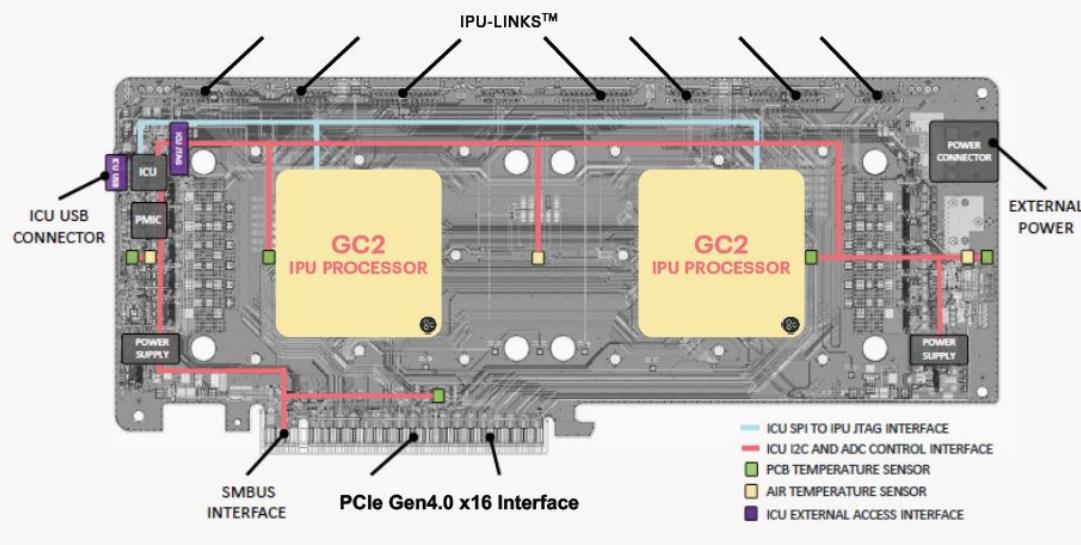
- Piotr Luszczek:

- Performance engineer and developer for multiple numerical libraries and benchmarks;
 - Member of the xSDK4ECP project;



Modern Hardware: Beyond Traditional IEEE Floating point

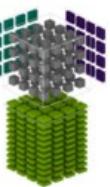
C2 IPU-PROCESSOR PCIe CARD



Floating point formats and performance on GPUs

2006	2007	2008	2009	2010	2011	2012	2013	2014	2015	2016	2017	2018	2019	2020
NVIDIA GPU generation	Tesla	Fermi		Kepler		Maxwell		Pascal		Volta				
Rel. compute performance	1 : 8	1 : 8		1 : 24		1 : 32		1 : 2 : 4		1 : 2 : 16*				double : single : half
Rel. memory performance	1 : 2	1 : 2		1 : 2		1 : 2		1 : 2 : 4		1 : 2 : 4				

*Tensor cores

$$D = \begin{pmatrix} A_{0,0} & A_{0,1} & A_{0,2} & A_{0,3} \\ A_{1,0} & A_{1,1} & A_{1,2} & A_{1,3} \\ A_{2,0} & A_{2,1} & A_{2,2} & A_{2,3} \\ A_{3,0} & A_{3,1} & A_{3,2} & A_{3,3} \end{pmatrix}_{\substack{\text{HMMA} \\ \text{IMMA}}} \begin{pmatrix} B_{0,0} & B_{0,1} & B_{0,2} & B_{0,3} \\ B_{1,0} & B_{1,1} & B_{1,2} & B_{1,3} \\ B_{2,0} & B_{2,1} & B_{2,2} & B_{2,3} \\ B_{3,0} & B_{3,1} & B_{3,2} & B_{3,3} \end{pmatrix}_{\substack{\text{FP16 or FP32} \\ \text{INT32}}} + \begin{pmatrix} C_{0,0} & C_{0,1} & C_{0,2} & C_{0,3} \\ C_{1,0} & C_{1,1} & C_{1,2} & C_{1,3} \\ C_{2,0} & C_{2,1} & C_{2,2} & C_{2,3} \\ C_{3,0} & C_{3,1} & C_{3,2} & C_{3,3} \end{pmatrix}_{\substack{\text{FP16 or FP32} \\ \text{INT32}}} \quad \begin{array}{l} \text{FP16} \\ \text{INT8 or UINT8} \end{array}$$


Floating point formats and performance on GPUs

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Rel. memory performance	1 : 2	1 : 2		1 : 2		1 : 2		1 : 2 : 4		1 : 2 : 4				

For **compute-bound applications**, the performance gains from using lower precision **depend on the architecture**.

Up to 16x for FP16 on Volta, up to 32x for FP32 on Maxwell.

For **memory-bound applications**, the performance gains from using lower precision are **architecture-independent** and correspond to the floating point format complexity (#bits).

Generally, 2x for FP32, 4x for FP16.

*Tensor cores

$$D = \begin{pmatrix} A_{0,0} & A_{0,1} & A_{0,2} & A_{0,3} \\ A_{1,0} & A_{1,1} & A_{1,2} & A_{1,3} \\ A_{2,0} & A_{2,1} & A_{2,2} & A_{2,3} \\ A_{3,0} & A_{3,1} & A_{3,2} & A_{3,3} \end{pmatrix} \begin{pmatrix} B_{0,0} & B_{0,1} & B_{0,2} & B_{0,3} \\ B_{1,0} & B_{1,1} & B_{1,2} & B_{1,3} \\ B_{2,0} & B_{2,1} & B_{2,2} & B_{2,3} \\ B_{3,0} & B_{3,1} & B_{3,2} & B_{3,3} \end{pmatrix} + \begin{pmatrix} C_{0,0} & C_{0,1} & C_{0,2} & C_{0,3} \\ C_{1,0} & C_{1,1} & C_{1,2} & C_{1,3} \\ C_{2,0} & C_{2,1} & C_{2,2} & C_{2,3} \\ C_{3,0} & C_{3,1} & C_{3,2} & C_{3,3} \end{pmatrix}$$

FP16 or FP32
INT32

FP16 or FP32
INT8 or UINT8

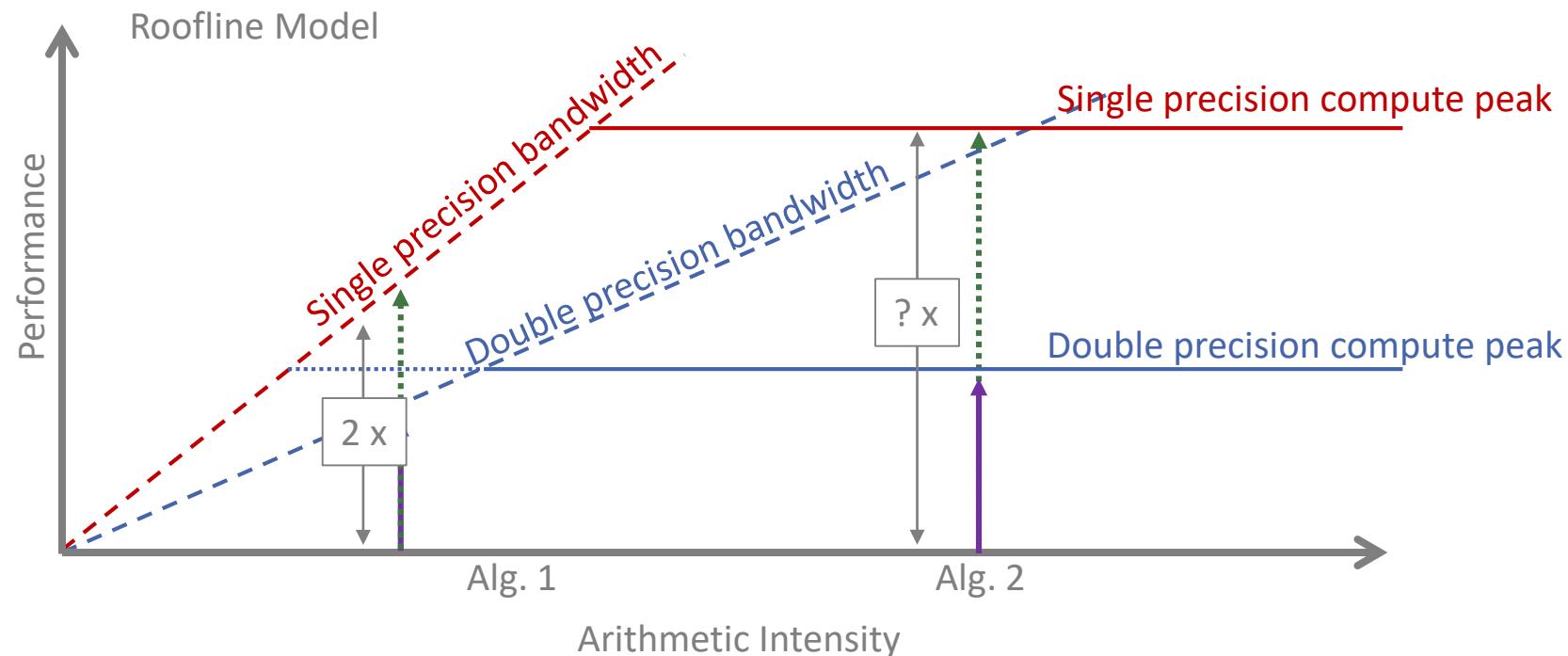
FP16 or FP32
INT8 or UINT8

FP16 or FP32
INT32



Take-Away I

- Performance of **compute-bound** algorithms depends on **format support of hardware**.
- Performance of **memory-bound** algorithms scales **hardware-independent** with **inverse of format complexity**.



IEEE 754 Floating Point Formats



Broadly speaking....

- The length of the **exponent** determines the **range** of the values that can be represented;
- The length of the **significand** determines how **accurate** values can be represented;

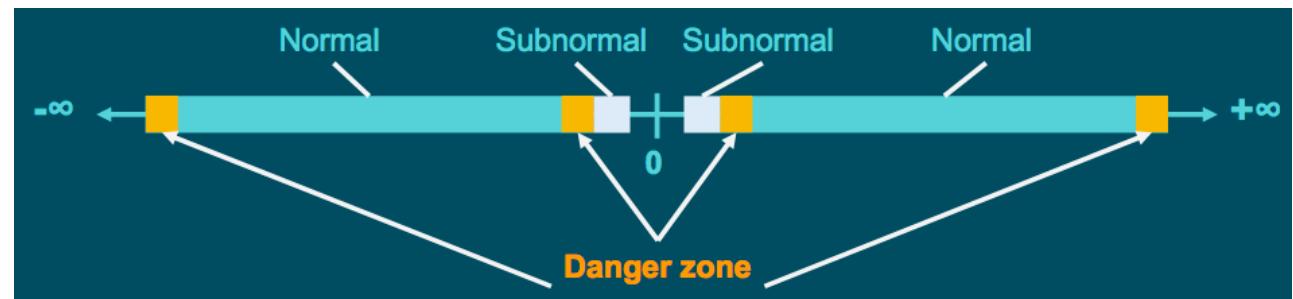
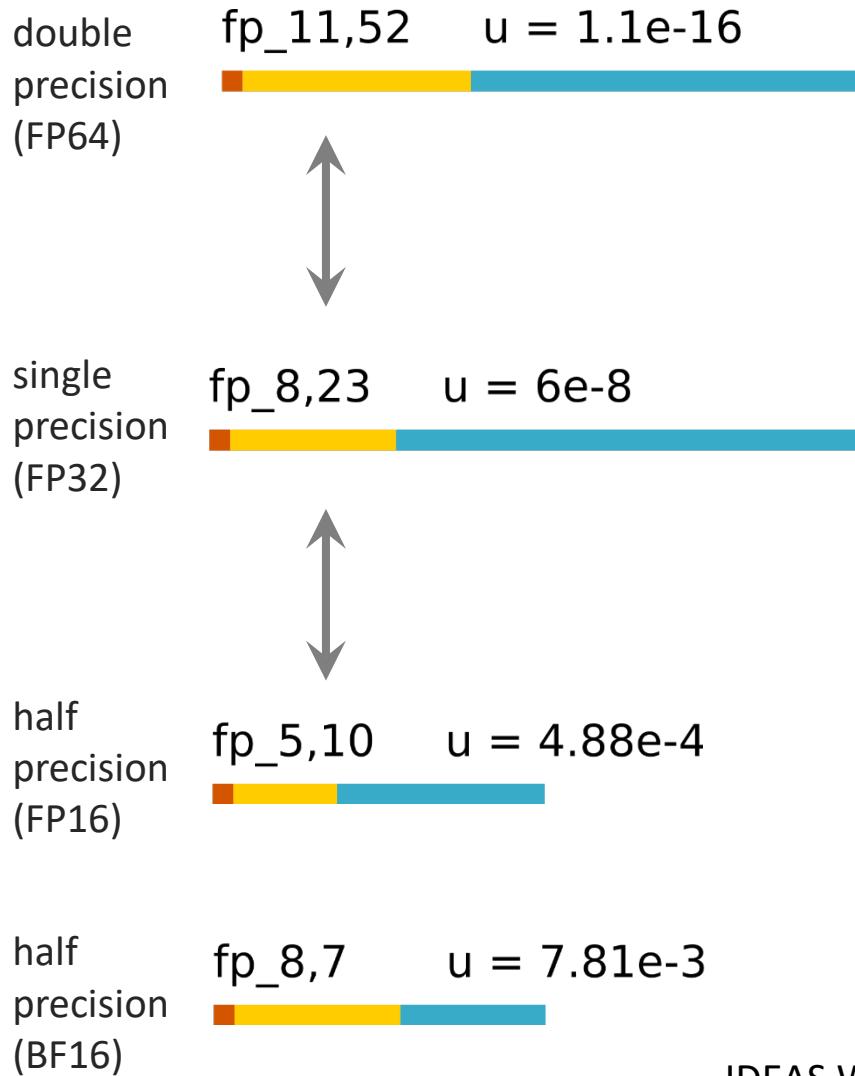


Figure courtesy of Ignacio Laguna, LLNL

IDEAS Webinar #34 by Ignacio Laguna on *Tools and Techniques for Floating-Point Analysis*

IEEE 754 Floating Point Formats



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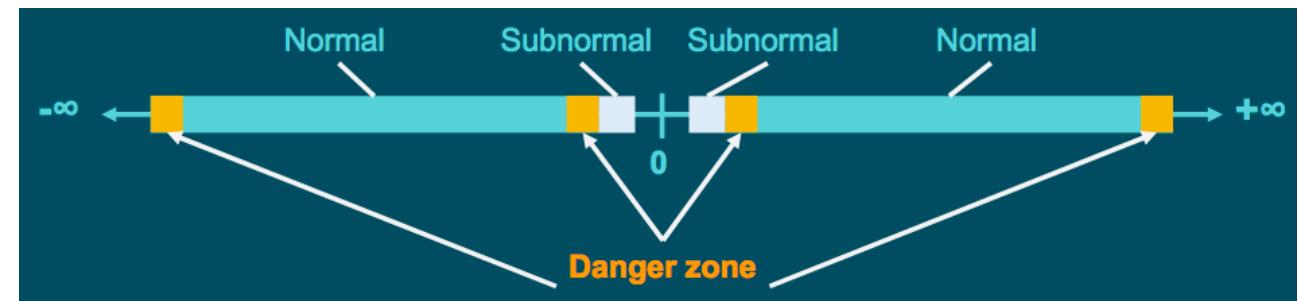


Figure courtesy of Ignacio Laguna, LLNL

IDEAS Webinar #34 by Ignacio Laguna on *Tools and Techniques for Floating-Point Analysis*

Floating point formats and accuracy

- The length of the **exponent** determines the **range** of the values that can be represented;
- The length of the **significand** determines how **accurate** values can be represented;
- *Rounding effects accumulate over a sequence of computations;*

Worst case: $fl \left(\sum_{i=1}^n x_i \right) = \sum_{i=1}^n x_i + n \cdot u$ with u being the unit round off.

N. Higham: Accuracy and stability of numerical algorithms. SIAM, 2002.

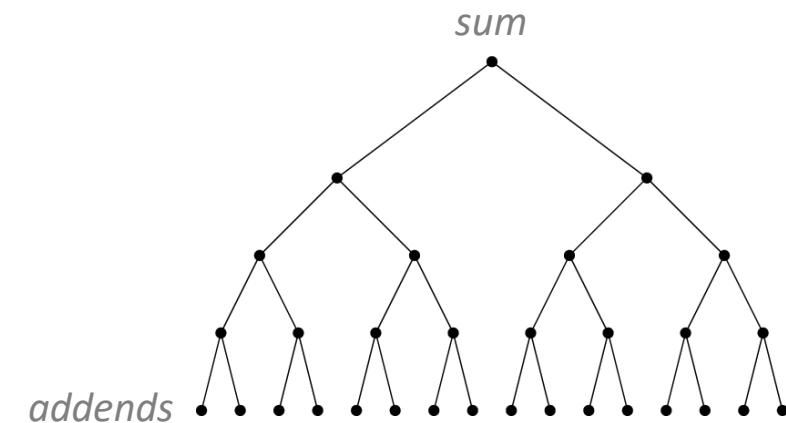
Floating point formats and accuracy

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Worst case: $fl \left(\sum_{i=1}^n x_i \right) = \sum_{i=1}^n x_i + n \cdot u$ with u being the unit round off.

In reality:

- *Stochastic effects reduce the impact of rounding;*
- *Parallel systems compute sums using blocking techniques (tree-based sum computation);*



N. Higham: Accuracy and stability of numerical algorithms. SIAM, 2002.

Floating point formats and accuracy

- The length of the **exponent** determines the **range** of the values that can be represented;
- The length of the **significand** determines how **accurate** values can be represented;
- *Rounding effects accumulate over a sequence of computations;*

Let us focus on linear systems of the form $Ax = b$.

- The conditioning of a linear system reflects how sensitive the solution x is with regard to changes in the right-hand side b .
- Rounding in the arithmetic operations of a linear solver equivalent to perturbations of the right-hand-side.
- Rule of thumb:

$$\text{relative residual accuracy} = (\text{unit round-off}) * (\text{linear system's condition number})$$

$$10^{-6} = 10^{-16} * 10^{10}$$

N. Higham: Accuracy and stability of numerical algorithms. SIAM, 2002.

Floating point formats and accuracy

Linear System $Ax=b$ with $\text{cond}(A) \approx 10^4$

Experiments based on the Ginkgo library <https://ginkgo-project.github.io/ginkgo/examples/adaptiveprecision-blockjacobi/adaptiveprecision-blockjacobi.cpp>

Floating point formats and accuracy

Linear System $Ax=b$ with $\text{cond}(A) \approx 10^4$

Double Precision

```
Initial residual norm sqrt(r^T r):
%%MatrixMarket matrix array real general
1 1
111.127
Final residual norm sqrt(r^T r):
%%MatrixMarket matrix array real general
1 1
5.0775e-10      Rel. Residual ~10^-12
CG iteration count:    1231
CG execution time [ms]: 140.038
```

relative residual accuracy = (unit round-off) * (linear system's condition number)



Experiments based on the Ginkgo library <https://ginkgo-project.github.io/ginkgo/examples/adaptiveprecision-blockjacobi/adaptiveprecision-blockjacobi.cpp>

Floating point formats and accuracy

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1 1
5.0775e-10    Rel. Residual ~10^-12
CG iteration count: 1231
CG execution time [ms]: 140.038
```

Exploring floating point formats in sparse iterative solvers:



<https://github.com/ginkgo-project/ginkgo>

- ValueType = double;
- + ValueType = float;

relative residual accuracy = (unit round-off) * (linear system's condition number)



Experiments based on the Ginkgo library <https://ginkgo-project.github.io/ginkgo/examples/adaptiveprecision-blockjacobi/adaptiveprecision-blockjacobi.cpp>

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CG execution time [ms]: 140.038
```

Single Precision

```
Initial residual norm sqrt(r^T r):  
%%MatrixMarket matrix array real general  
1 1  
111.127  
Final residual norm sqrt(r^T r):  
%%MatrixMarket matrix array real general  
1 1  
0.179829 Rel. Residual ~10^-4  
CG iteration count: 1234  
CG execution time [ms]: 127.152
```

relative residual accuracy = (unit round-off) * (linear system's condition number)



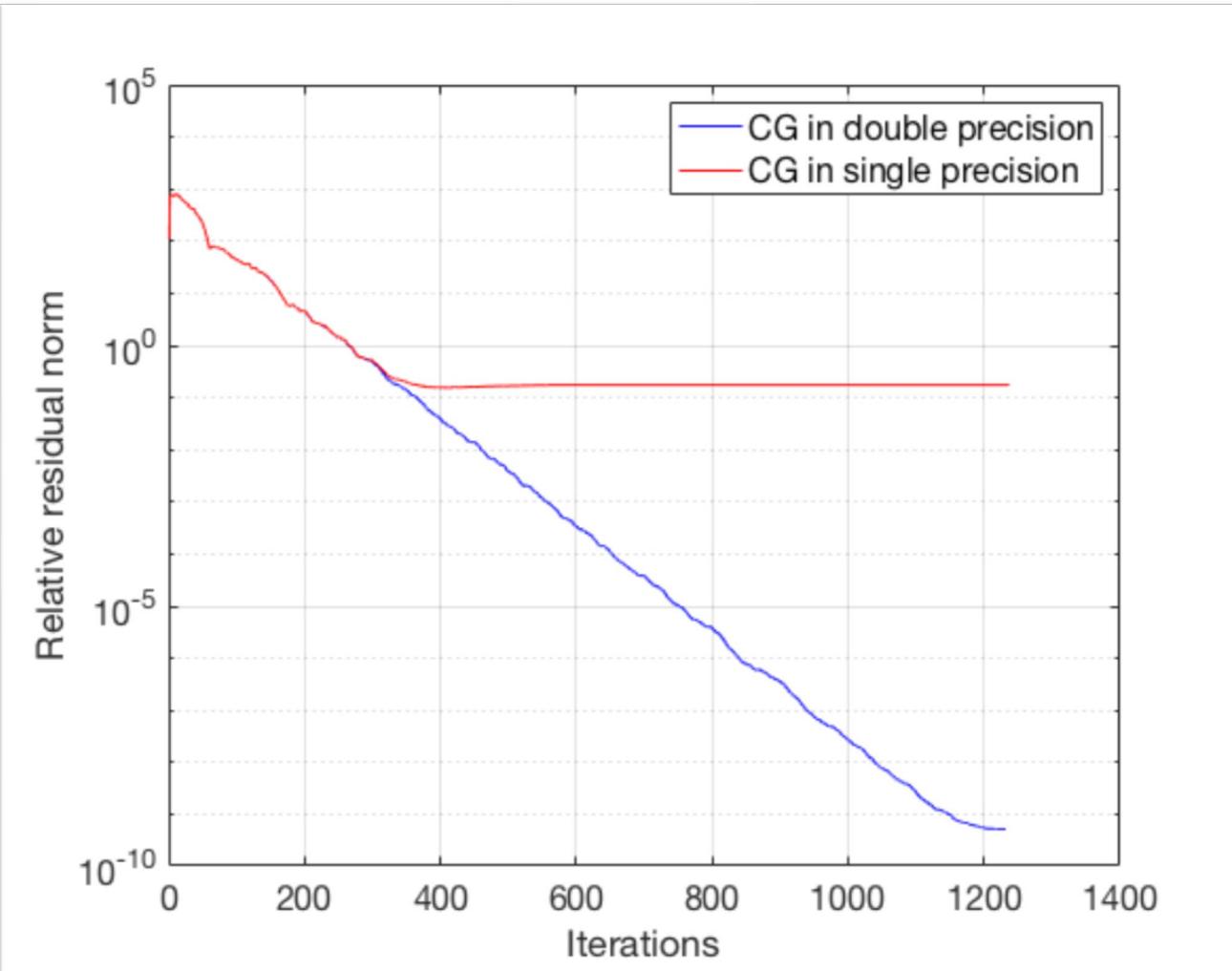
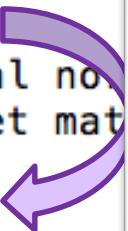
Experiments based on the Ginkgo library <https://ginkgo-project.github.io/ginkgo/examples/adaptiveprecision-blockjacobi/adaptiveprecision-blockjacobi.cpp>

Floating point formats and accuracy

Linear System $Ax=b$ with

Double Precision

```
Initial residual norm: 111.127
%%MatrixMarket matrix
1 1
Final residual norm: 5.0775e-10
%%MatrixMarket matrix
1 1
CG iteration count: 1234
CG execution time: 127.152
```



```
rt(r^T r):
ray real general
(r^T r):
ray real general
al ~10^-4
1234
127.152
```

Experiments based on the Ginkgo library <https://ginkgo-project.github.io/ginkgo/examples/adaptiveprecision-blockjacobi/adaptiveprecision-blockjacobi.cpp>

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1 1  
5.0775e-10  
CG iteration count: 1231  
CG execution time [ms]: 140.038
```

Single Precision

```
Initial residual norm sqrt(r^T r):  
%%MatrixMarket matrix array real general  
1 1  
111.127  
Final residual norm sqrt(r^T r):  
%%MatrixMarket matrix array real general  
1 1  
0.179829  
CG iteration count: 1234  
CG execution time [ms]: 127.152
```

Single Precision is 10% faster!

Experiments based on the Ginkgo library <https://ginkgo-project.github.io/ginkgo/examples/adaptiveprecision-blockjacobi/adaptiveprecision-blockjacobi.cpp>

Floating point formats and accuracy

Linear System $Ax=b$ with $\text{cond}(A) \approx 10^7$

apache2 from SuiteSparse

Double Precision

```
Initial residual norm sqrt(r^T r):  
%%MatrixMarket matrix array real general  
1 1  
1390.67  
Final residual norm sqrt(r^T r):  
%%MatrixMarket matrix array real general  
1 1  
4.53915e-06 Rel. Residual ~10-9  
CG iteration count: 6460  
CG execution time [ms]: 2992.91
```

Single Precision

```
Initial residual norm sqrt(r^T r):  
%%MatrixMarket matrix array real general  
1 1  
1390.67  
Final residual norm sqrt(r^T r):  
%%MatrixMarket matrix array real general  
1 1  
1588.77 No improvement  
CG iteration count: 8887  
CG execution time [ms]: 2972.46
```

Experiments based on the Ginkgo library <https://ginkgo-project.github.io/>

ginkgo/examples/adaptiveprecision-blockjacobi/adaptiveprecision-blockjacobi.cpp

Take-Away II

- Relative residual accuracy = (unit round-off) * (linear system's condition number)
- For ill-conditioned problems, we need high precision to provide high accuracy results.
- Only if the problem is well-conditioned, and a low-accuracy solution is acceptable, we can use a low precision format throughout the complete solution process.
- Templating the precision format (i.e. ValueType) allows to quickly switch between formats.
 - C++ very powerful in this respect
 - Use production-ready libraries templating the precision: Ginkgo, Kokkos Kernels, Trilinos, etc.

Low precision for solving ill-conditioned problems

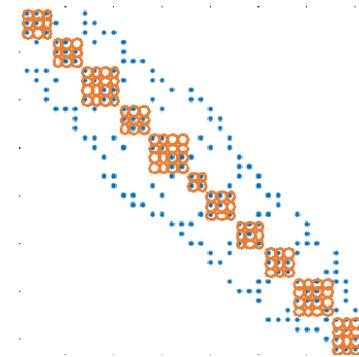
- **Preconditioning iterative solvers.**
 - Idea: Approximate inverse of system matrix to make the system “easier to solve”: $P^{-1} \approx A^{-1}$ and solve $Ax = b \Leftrightarrow P^{-1}Ax = P^{-1}b \Leftrightarrow \tilde{A}x = \tilde{b}$.

Low precision for solving ill-conditioned problems

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 - Why should we use a preconditioner P^{-1} in full (high) precision?

Low precision for solving ill-conditioned problems

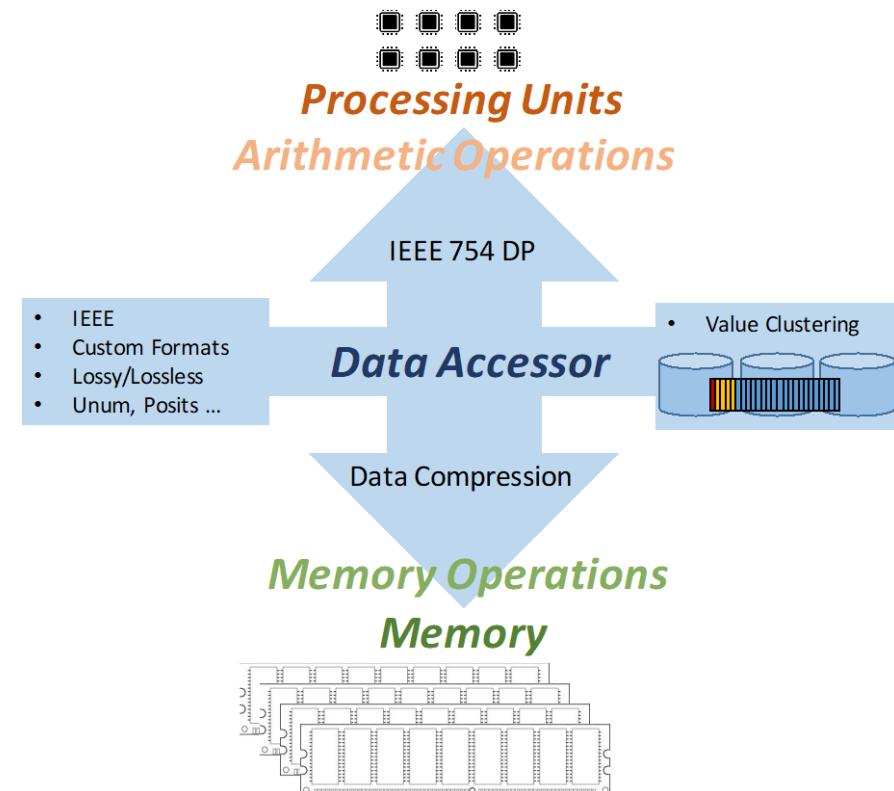
- Preconditioning iterative solvers.
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 - Why should we use a preconditioner P^{-1} in full (high) precision?
- Jacobi method based on diagonal scaling $P = \text{diag}(A)$
- Block-Jacobi is based on block-diagonal scaling: $P = \text{diag}_B(A)$
 - Each block corresponds to one (small) linear system.
 - Larger blocks typically improve convergence.
 - Larger blocks make block-Jacobi more expensive.



Low precision for solving ill-conditioned problems

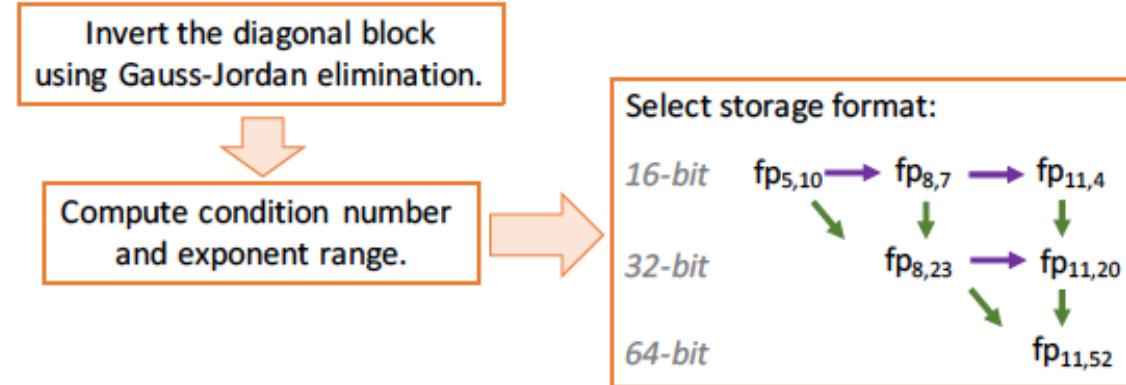
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 - Larger blocks make block-Jacobi more expensive.

Idea: Store the inverted diagonal in low precision



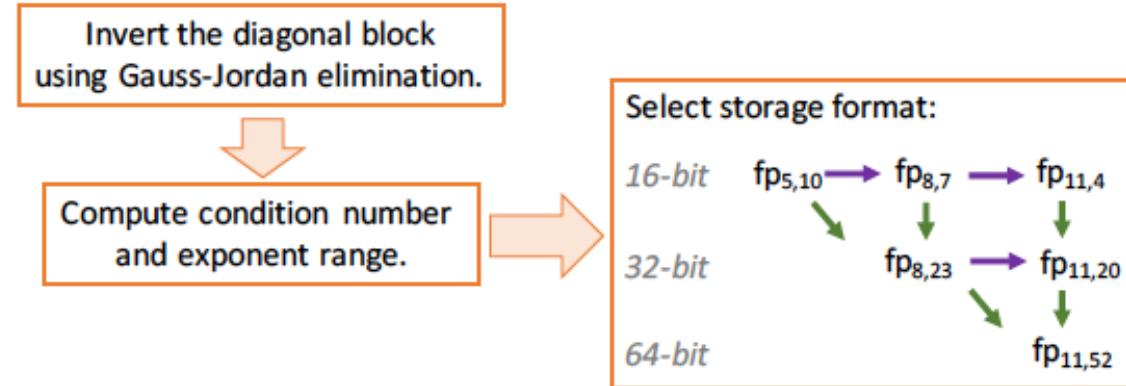
Adaptive Precision Preconditioning

- Choose how much accuracy of the preconditioner should be preserved by the storage format.
- All computations use double precision, but store blocks in lower precision.



Adaptive Precision Preconditioning

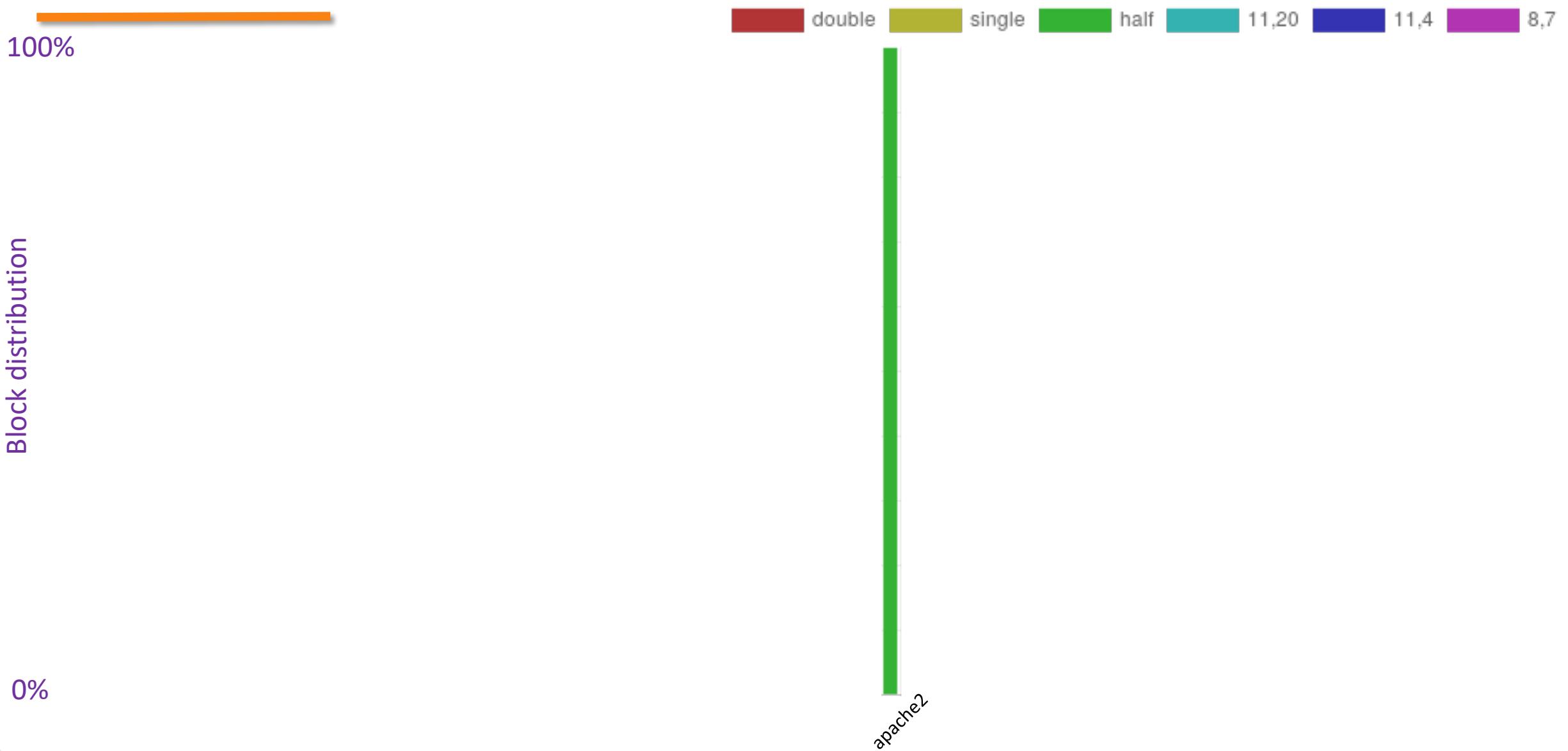
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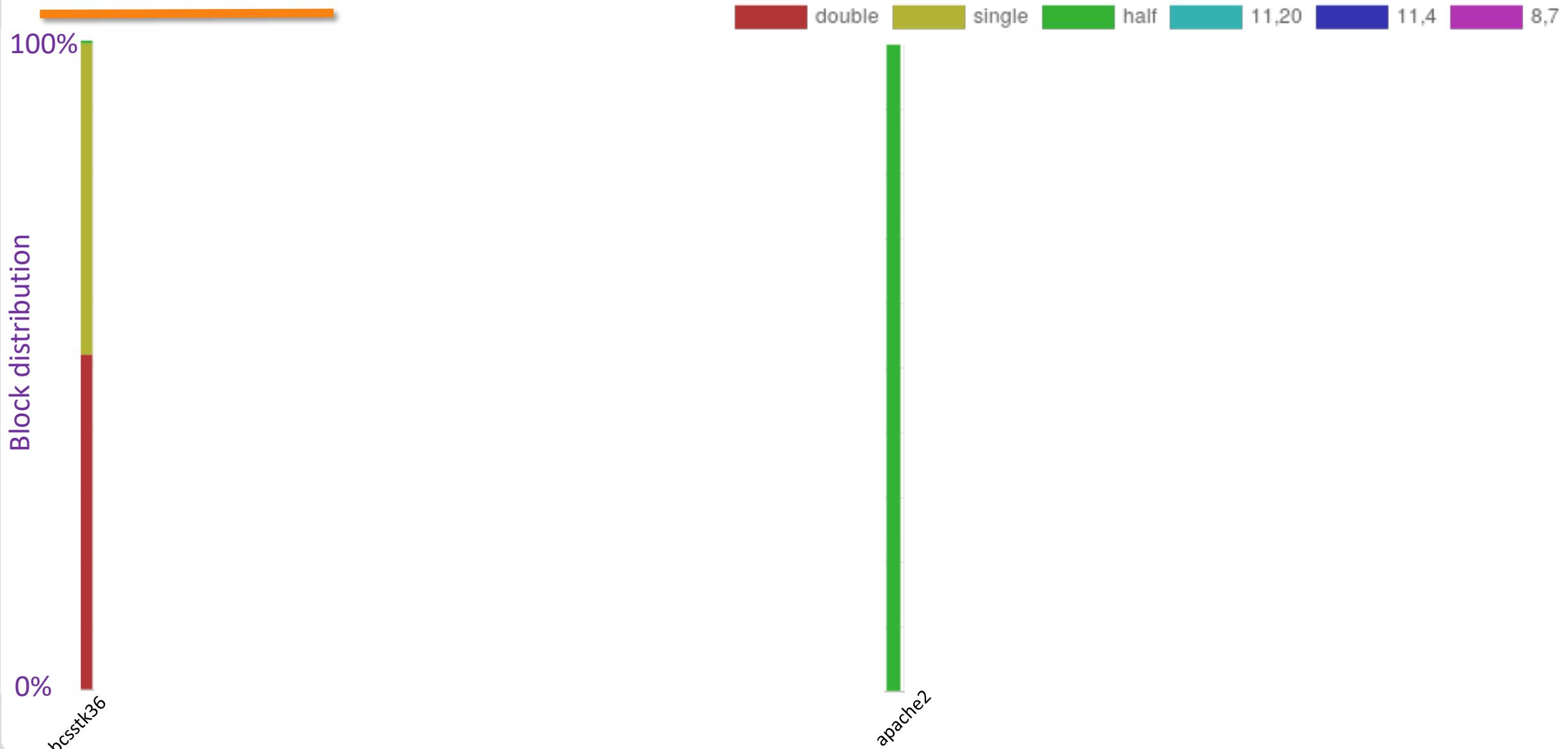
- + Regularity preserved;
- + Flexible in the accuracy preserved;
Preserving Preconditioner accuracy of 10⁻²
- + No flexible Krylov solver needed
(Preconditioner constant operator);
- + Can handle non-spd problems (featuring pivoting);
- + Can be used in any preconditionable solver;

- Overhead of the precision detection
(condition number calculation);
- Overhead from storing precision information
(need to additionally store/retrieve flag);
- Speedups / preconditioner quality problem-dependent;

Adaptive Precision Preconditioning



Adaptive Precision Preconditioning



Adaptive Precision Preconditioning



Floating point formats and accuracy

Linear System $Ax=b$ with $\text{cond}(A) \approx 10^7$

apache2 from SuiteSparse

Double Precision + Double Preconditioner

```
Initial residual norm sqrt(r^T r):  
%%MatrixMarket matrix array real general  
1 1  
1390.67  
Final residual norm sqrt(r^T r):  
%%MatrixMarket matrix array real general  
1 1  
3.97985e-06 Rel. Residual ~10^-9  
CG iteration count: 4797  
CG execution time [ms]: 2971.18
```

Double Precision + Mixed Precision Preconditioner

```
Initial residual norm sqrt(r^T r):  
%%MatrixMarket matrix array real general  
1 1  
1390.67  
Final residual norm sqrt(r^T r):  
%%MatrixMarket matrix array real general  
1 1  
3.98574e-06 Rel. Residual ~10^-9  
CG iteration count: 4794  
CG execution time [ms]: 2568.1
```

16% runtime improvement

Experiments based on the Ginkgo library <https://ginkgo-project.github.io/>

ginkgo/examples/adaptiveprecision-blockjacobi/adaptiveprecision-blockjacobi.cpp

Floating point formats and accuracy

ginkgo/examples/adaptiveprecision-blockjacobi/adaptiveprecision-blockjacobi.cpp

...

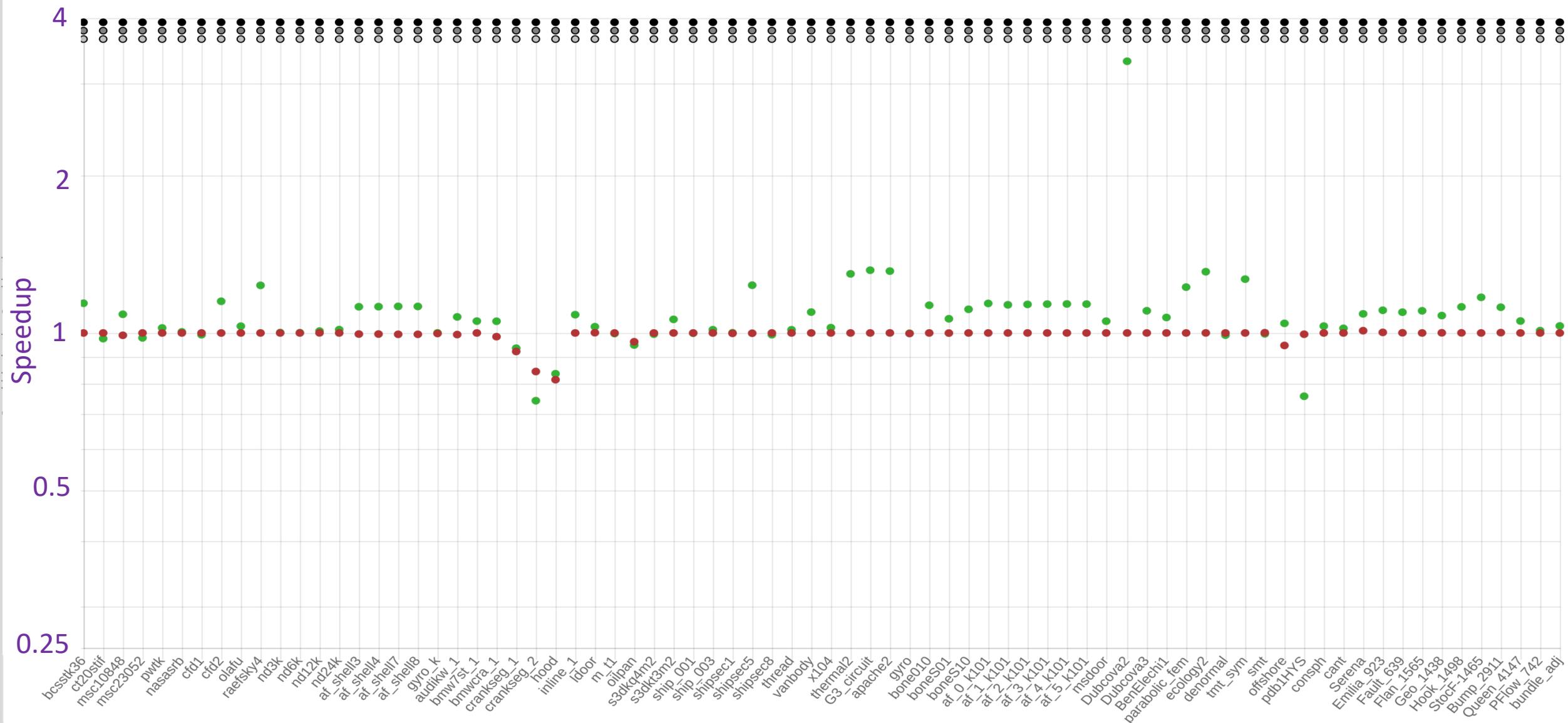
```
auto solver_gen =
    cg::build()
    .with_criteria(gko::share(iter_stop), gko::share(tol_stop))
        .with_preconditioner(bj::build())
    .with_max_block_size(16u)
    .with_storage_optimization(
        gko::precision_reduction::autodetect())
    .on(exec))
    .on(exec);
```

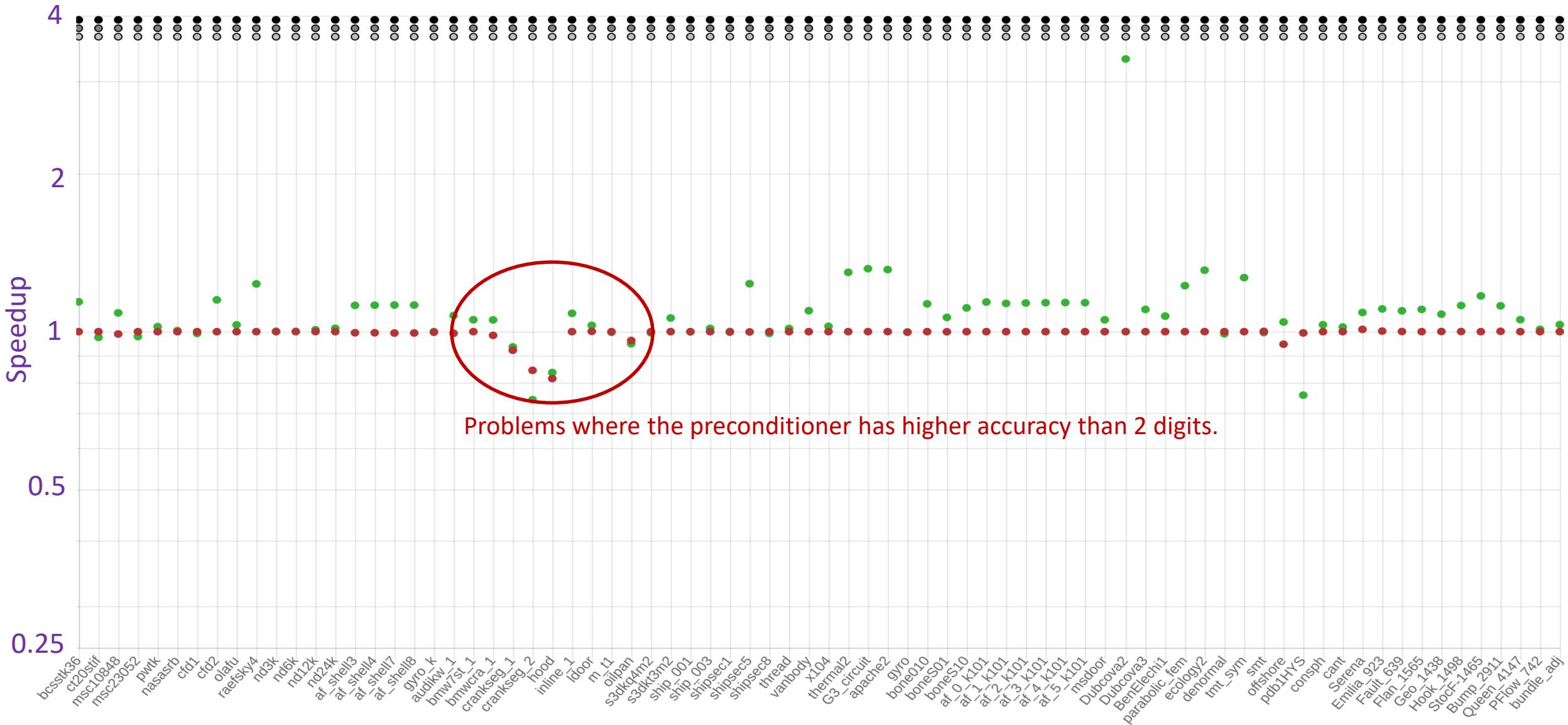
...

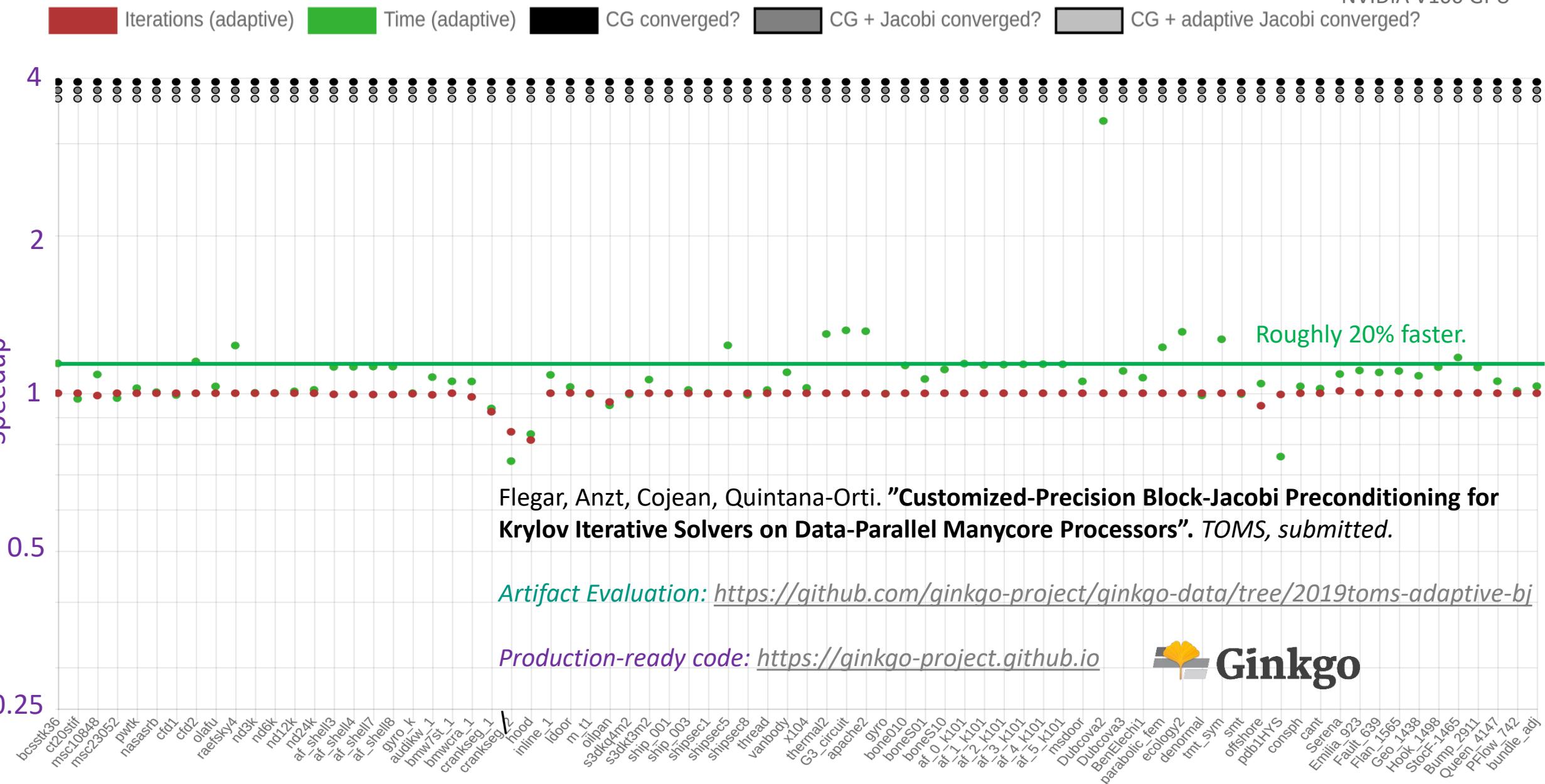
Experiments based on the Ginkgo library <https://ginkgo-project.github.io/>

[ginkgo/examples/adaptiveprecision-blockjacobi/adaptiveprecision-blockjacobi.cpp](https://ginkgo-project.github.io/ginkgo/examples/adaptiveprecision-blockjacobi/adaptiveprecision-blockjacobi.cpp)

Iterations (adaptive) Time (adaptive) CG converged? CG + Jacobi converged? CG + adaptive Jacobi converged?







Take-Away III

- Low precision preconditioners can be used to accelerate iterative solvers.
 - Preconditioners need to adapt their precision to numerical requirements.
 - The preconditioner precision determines how much accuracy is preserved.
- For memory-bound preconditioners, decoupling the arithmetic precision from the memory precision provides the runtime savings while preserving a constant preconditioner.
- To increase the performance benefits, shift most of the work to the low precision preconditioner.

Using a low precision solver as preconditioner

- To increase the performance benefits, shift most of the work to the low precision preconditioner.
- Use a simple (cheap) iterative solver in high precision and a sophisticated (expensive) solver in low precision as preconditioner.
 - Most of the work is done in low precision (fast).
 - The high precision outer solver ensures high quality of the solution.
- Popular example: Iterative Refinement

For an approximate solution $x^{(k)}$, the residual computes as $r = b - Ax^{(k)}$.
The exact solution for $Ax = b$ is $x = x^{(k)} + c$ where c is the solution of $Ac = r$.

Mixed Precision Iterative Refinement

```
Choose initial guess x      high precision
do {
    Compute r = b - Ax    high precision
    Solve A * c = r       low precision
    Update x = x + c     high precision
} while ( ||r|| > tol )   high precision
```

N. Higham: Accuracy and stability of numerical algorithms. SIAM, 2002.



Sri Pranesh's mixed precision Matlab suite:
https://github.com/SrikaraPranesh/Multi_precision_NLA_kernels

Mixed Precision Iterative Refinement using sparse iterative solvers

Linear System $Ax=b$ with $\text{cond}(A) \approx 10^4$

Double Precision Iterative Refinement

```
Initial residual norm sqrt(r^T r):  
%MatrixMarket matrix array real general  
1 1  
111.127  
Final residual norm sqrt(r^T r):  
%MatrixMarket matrix array real general  
1 1  
7.16102e-11 Rel. Residual ~10^-14  
MPIR iteration count: 18  
MPIR execution time [ms]: 213.491
```

Mixed Precision Iterative Refinement

```
Initial residual norm sqrt(r^T r):  
%MatrixMarket matrix array real general  
1 1  
111.127  
Final residual norm sqrt(r^T r):  
%MatrixMarket matrix array real general  
1 1  
7.41333e-11 Rel. Residual ~10^-14  
MPIR iteration count: 18  
MPIR execution time [ms]: 183.296
```

16% runtime improvement

Experiments based on the Ginkgo library <https://ginkgo-project.github.io/ginkgo/examples/mixed-precision-ir/mixed-precision-ir.cpp>

Mixed Precision Iterative Refinement using sparse iterative solvers

Some references:

Strzodka et al. Pipelined Mixed Precision Algorithms on FPGAs for Fast and Accurate PDE Solvers from Low Precision Components, IEEE Symposium on Field-Programmable Custom Computing Machines, 2006.

Goedekke et al. Performance and accuracy of hardware-oriented native-, emulated- and mixed-precision solvers in FEM simulations, International Journal of Parallel, Emergent and Distributed Systems, 2007.

Buratti et al. Using Mixed Precision for Sparse Matrix Computations to Enhance the Performance while Achieving 64-bit Accuracy, ACM TOMS 2008.

Baboulin et al. Accelerating scientific computations with mixed precision algorithms, CPC, 2009.

Anzt et al. Mixed precision iterative refinement methods for linear systems: Convergence analysis based on Krylov subspace methods, PARA 2010.

...

For sparse iterative methods, the benefits relate to the bandwidth savings.

Mixed Precision Iterative Refinement using sparse iterative solvers

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Strzodka et al. Pipelined Mixed Precision Algorithms on FPGAs for Fast and Accurate PDE Solvers from Low Precision Components, IEEE Symposium on Field-Programmable Custom Computing Machines, 2006.

Goedekke et al. Performance and accuracy of hardware-oriented native-, emulated- and mixed-precision solvers in FEM simulations, International Journal of Parallel, Emergent and Distributed Systems, 2007.

Buratti et al. Using Mixed Precision for Sparse Matrix Computations to Enhance the Performance while Achieving 64-bit Accuracy, ACM TOMS 2008.

Baboulin et al. Accelerating scientific computations with mixed precision algorithms, CPC, 2009.

Anzt et al. Mixed precision iterative refinement methods for linear systems: Convergence analysis based on Krylov subspace methods, PARA 2010.

...
For sparse iterative methods, the benefits relate to the bandwidth savings.

Some software packages allowing for mixed precision sparse linear system solves:



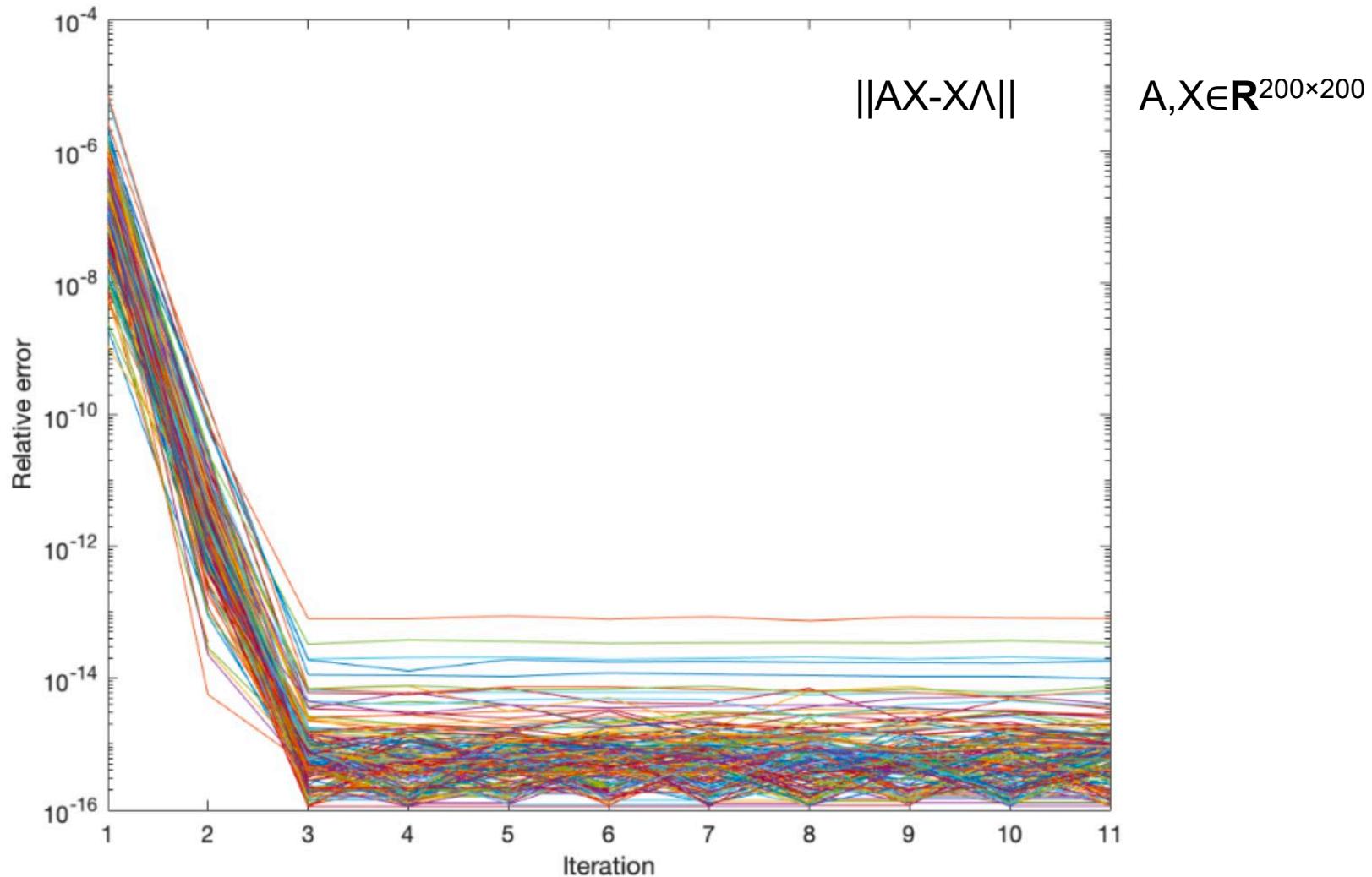
Refinement symmetric eigenvalues and eigenvectors

```

1: Input:  $A = A^T \in \mathbb{R}^{n \times n}$ ,  $\widehat{X} \in \mathbb{R}^{n \times \ell}$ ,  $1 \leq \ell \leq n$  ◀ Low precision inputs (must be accurate in lower precision)
2: Output:  $X' \in \mathbb{R}^{n \times \ell}$ ,  $\widetilde{D} = \text{diag}(\widetilde{\lambda}_i) \in \mathbb{R}^{\ell \times \ell}$ ,  $\widetilde{E} \in \mathbb{R}^{\ell \times \ell}$ ,  $\omega \in \mathbb{R}$ 
3: function  $[X', \widetilde{D}, \widetilde{E}, \omega] \leftarrow \text{REFSYEv}(A, \widehat{X})$ 
4:    $R \leftarrow \mathbb{I}_n - \widehat{X}^T \widehat{X}$ 
5:    $S \leftarrow \widehat{X}^T A \widehat{X}$ 
6:    $\widetilde{\lambda}_i \leftarrow s_{ii}/(1 - r_{ii})$  for  $i = 1, \dots, \ell$  ▷ Compute approximate eigenvalues.
7:    $\widetilde{D} \leftarrow \text{diag}(\widetilde{\lambda}_i)$ 
8:    $\omega \leftarrow 2 \left( \|S - \widetilde{D}\|_2 + \|A\|_2 \|R\|_2 \right)$ 
9:    $e_{ij} \leftarrow \begin{cases} \frac{s_{ij} + \widetilde{\lambda}_j r_{ij}}{\widetilde{\lambda}_j - \widetilde{\lambda}_i} & \text{if } |\widetilde{\lambda}_i - \widetilde{\lambda}_j| > \omega \\ r_{ij}/2 & \text{otherwise} \end{cases}$  for  $1 \leq i, j \leq \ell$  ▷ Compute the entries of the refinement matrix  $\widetilde{E}$ .
10:   $X' \leftarrow \widehat{X} + \widehat{X} \widetilde{E}$  ▷ Update  $\widehat{X}$  by  $\widehat{X}(\mathbb{I}_n + \widetilde{E})$ 
11: end function

```

Convergence for each of 200 eigenvalues (backward error)



Complexity Analysis

```
1: Input:  $A = A^T \in \mathbb{R}^{n \times n}$ ,  $\widehat{X} \in \mathbb{R}^{n \times \ell}$ ,  $1 \leq \ell \leq n$ 
2: Output:  $X' \in \mathbb{R}^{n \times \ell}$ ,  $\widetilde{D} = \text{diag}(\widetilde{\lambda}_i) \in \mathbb{R}^{\ell \times \ell}$ ,  $\widetilde{E} \in \mathbb{R}^{\ell \times \ell}$ ,  $\omega \in \mathbb{R}$ 
3: function  $[X', \widetilde{D}, \widetilde{E}, \omega] \leftarrow \text{REFSYEv}(A, \widehat{X})$ 
4:    $R \leftarrow \mathbb{I}_n - \widehat{X}^T \widehat{X}$ 
5:    $S \leftarrow \widehat{X}^T A \widehat{X}$ 
6:    $\widetilde{\lambda}_i \leftarrow s_{ii}/(1 - r_{ii})$  for  $i = 1, \dots, \ell$ 
7:    $\widetilde{D} \leftarrow \text{diag}(\widetilde{\lambda}_i)$ 
8:    $\omega \leftarrow 2 \left( \|S - \widetilde{D}\|_2 + \|A\|_2 \|R\|_2 \right)$ 
9:    $e_{ij} \leftarrow \begin{cases} \frac{s_{ij} + \widetilde{\lambda}_j r_{ij}}{\widetilde{\lambda}_j - \widetilde{\lambda}_i} & \text{if } |\widetilde{\lambda}_i - \widetilde{\lambda}_j| > \omega \\ r_{ij}/2 & \text{otherwise} \end{cases}$  for  $1 \leq i, j \leq \ell$ 
10:   $X' \leftarrow \widehat{X} + \widetilde{X} \widetilde{E}$ 
11: end function
```

4 matrix multiplies in higher precision $7n^3$ per iteration

► Compute approximate eigenvalues.

► Compute the entries of the refinement matrix \widetilde{E} .

► Update \widehat{X} by $\widehat{X}(\mathbb{I}_n + \widetilde{E})$

Take-Away IV

- Mixed precision iterative refinement is a powerful strategy to accelerate linear solves;
 - Iterative inner solver e.g. for sparse systems;
 - Direct inner solver e.g. for dense systems;
- Mixed precision iterative refinement can also be used for eigenvalue problems;
 - Low precision eigenvector approximations as input;
 - Convergence in high precision after 3-4 IR steps;
- The performance benefits depend on the problem and hardware capabilities;

References and further reading

- Higham. **Accuracy and stability of numerical algorithms**, SIAM, 2002.
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